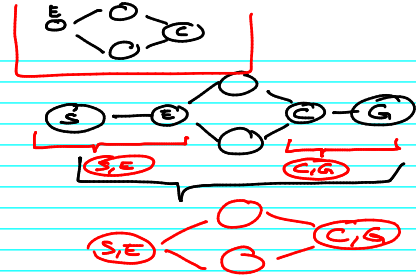


Map(f) →
pipeline(seq(split), f, seq(gather))

pipe(map(f), map(g)) →
map(pipeline(f, g))



DOUBLE/TRIPLE BUFFERING

→ COMM. OVERHEADS ↓

```

Stage i
→ ○ →
write (not termination) }
task = receive();
res = compute(task);
send (result);
}
    
```

$$\text{Latency} = T_{\text{comm}} + T_{\text{task}} + T_{\text{comm}}$$

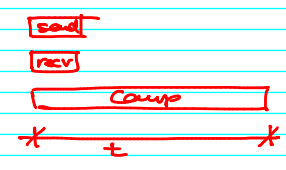
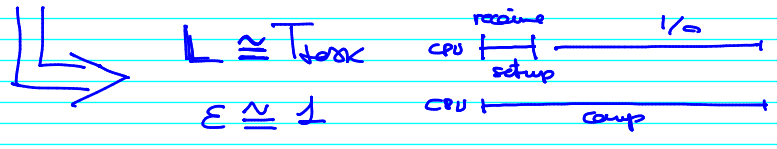
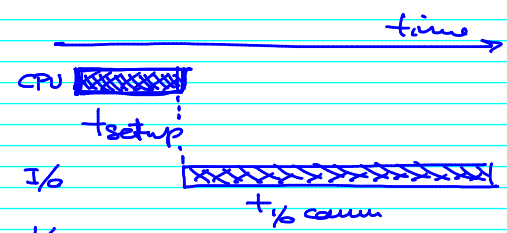
overhead (points to the first T_{comm})
useful computations (points to T_{task})

$$E = \frac{T_{\text{task}}}{2T_{\text{comm}} + T_{\text{task}}}$$

TRIPLE BUFF

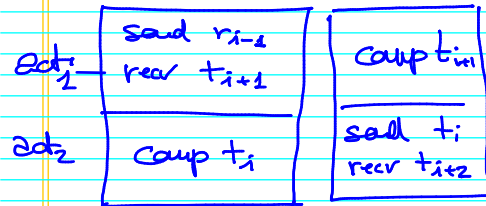
act ₁	receive	comp	send	reain
act ₂	comp	send	reain	comp
act ₃	send	reain	comp	send

← t t t t →
time



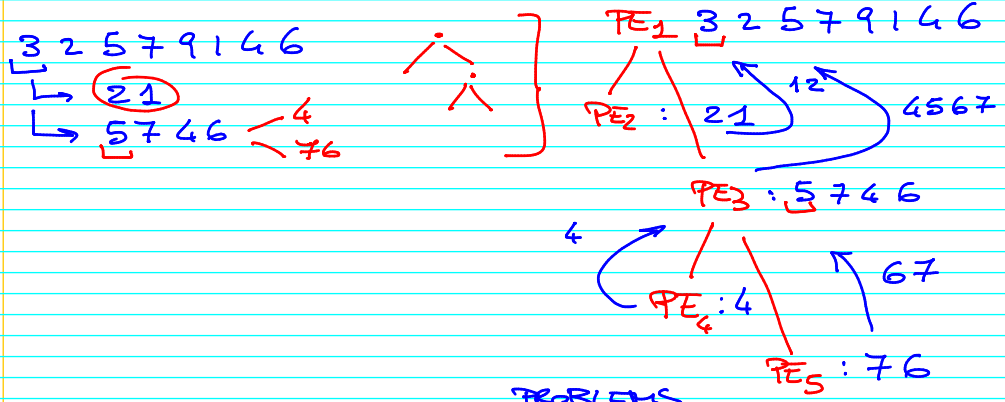
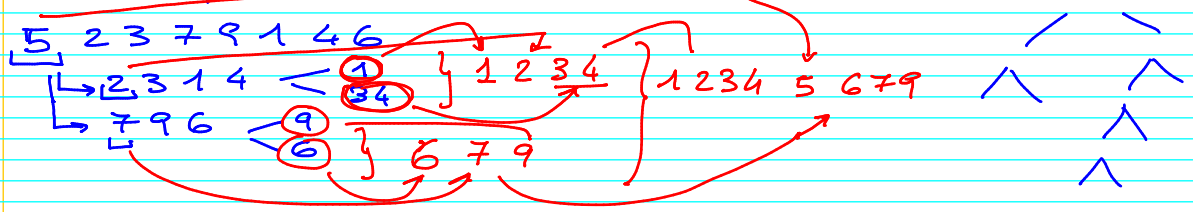
DOUBLE BUFFERING

↳ reduced version
of triple buffering



FARM WITH FEEDBACK

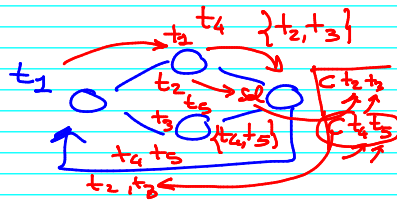
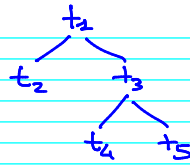
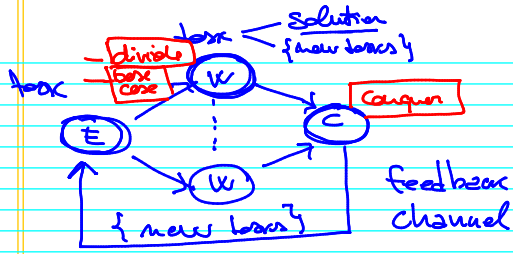
divide & conquer



PROBLEMS

- 1) RESOURCE RECRUITMENT
- 2) MANAGE FREE RESOURCE

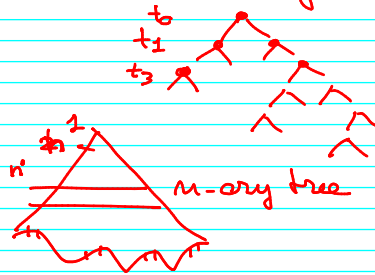
DYNAMIC RESOURCE MGT



```

div & conq bc bcs d c x =
  if (bc x)      (bcs x)
  else (conq
    map (d & c ... (split x));
  )
  
```

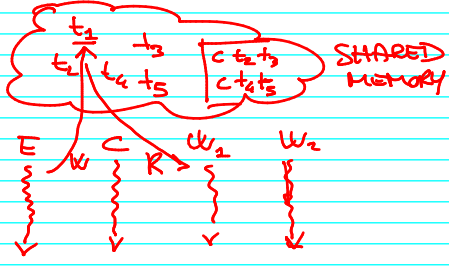
Performance



there is a max problem degree that depends on the computation AND on the input data

FASTFLOW

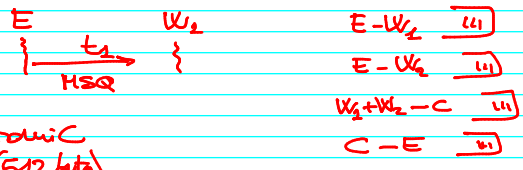
DLC (FARM WITH FEEDBACK)
ON A MULTICORE?



Message queues between "concurrent activities" (t_1 reads)

- ① only contain pointers
- ② guarantee synchronization

pipe writes & reads are guaranteed to be atomic iff $\dim(\text{data}) < \text{CONSTANT}$ (512 bytes at least)



```

E          W2
retcode = pipe (fd);
pthread_create < t1, t2
t1          t2
write (fd[2], ...) read (fd[0], ...)
    
```

```

#include <unistd.h>
retcode = pipe (int [2])
syscall

fd [ ] ← descriptor write a file
         ← " " read a file
W → [ ] → R
    
```