

Specie di pattern

concurrency
algorithms

Support structure

recursion

} skeleton

} template

building blocks

$$\{ \text{Building block} \} = \{ \text{Wrapper} \} \cup \{ \text{functor} \} \cup \{ \text{combinator} \}$$

WRAPPER

combinatori: codice \rightarrow "componente building block"

Wrapper di codice seq ((code)) ((f))

*) 1 ingresso e 1 uscita



*) statero

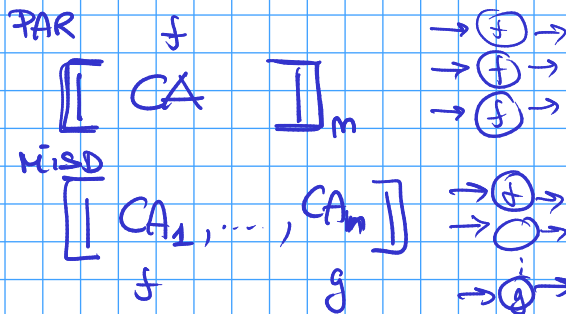
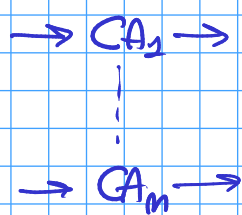
Wrapper di codice parallel ((f)) \rightarrow (f) \rightarrow { } { } { }

*) 1 ingresso e 1 uscita

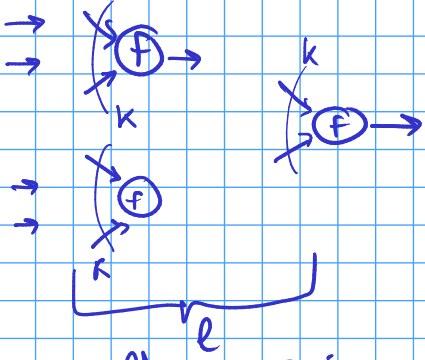
*) statero

\Rightarrow stato di codice

Funzionali (solo "quelli da calcolo")
(in parallelo)

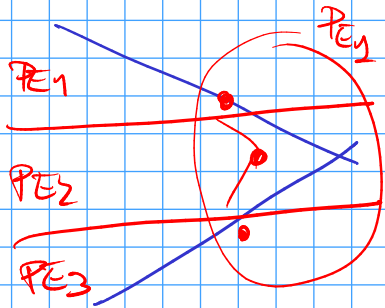
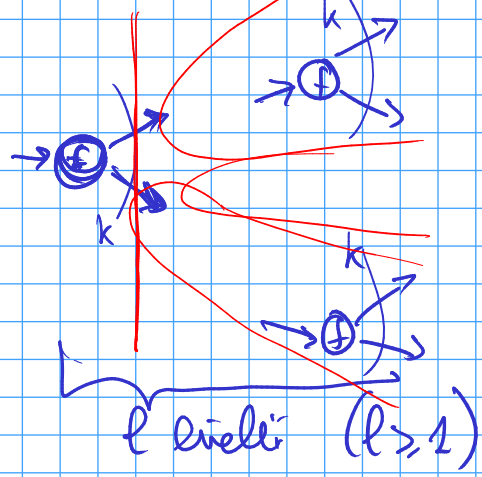


REDUCE $f \triangleleft$

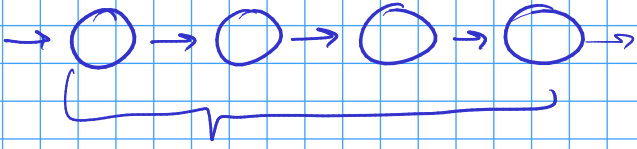


albero k -ario
profondità l ($l \geq 1$)

SPREAD $f \triangleright$



se f è commutativa e associativa



PIPELINE

COMBINATORI

1-to-N



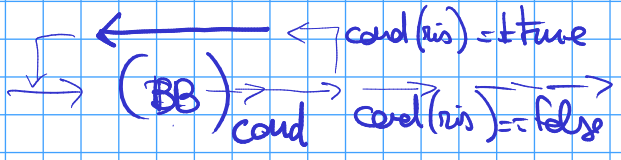
$Pole \in \{ \text{scatter, Broadcast, Unicast (P)} \}$
 $P = \text{auto, Rr, ...}$

N-to-1

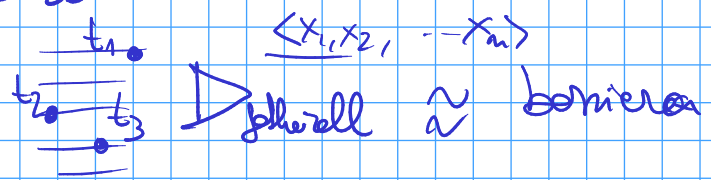


$Pole = \{ \text{gather, gather all, Reduce} \}$

feedback

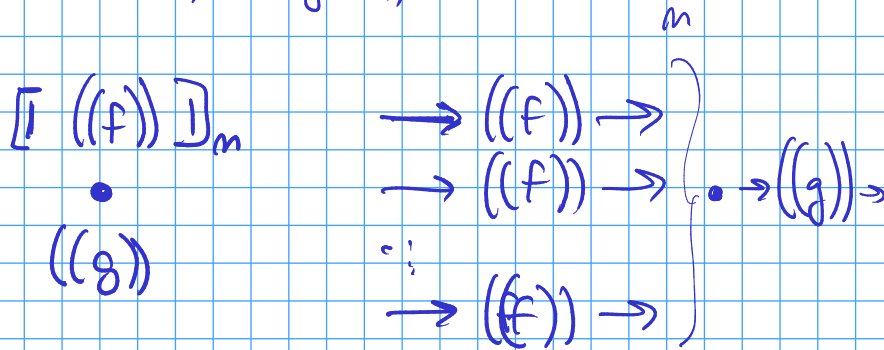
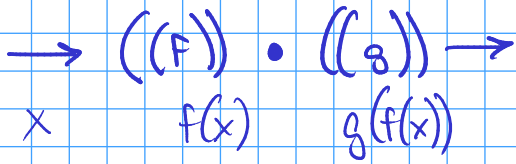


building block = BB



Composizionale

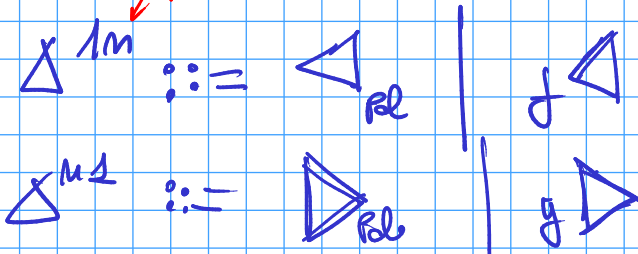
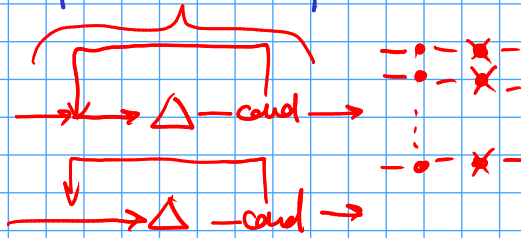
• ε pipeline (infinita)



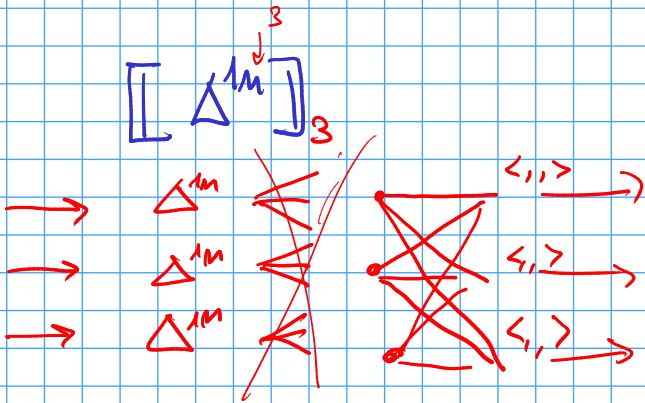
$$\Delta^m ::= [\Delta] \mid [\Delta_1 \dots \Delta_n] \mid \overleftarrow{\Delta^m_{\text{cand}}} \mid \Delta^m \cdot \Delta^m$$

per
quanti uguali
quante vacite

mid

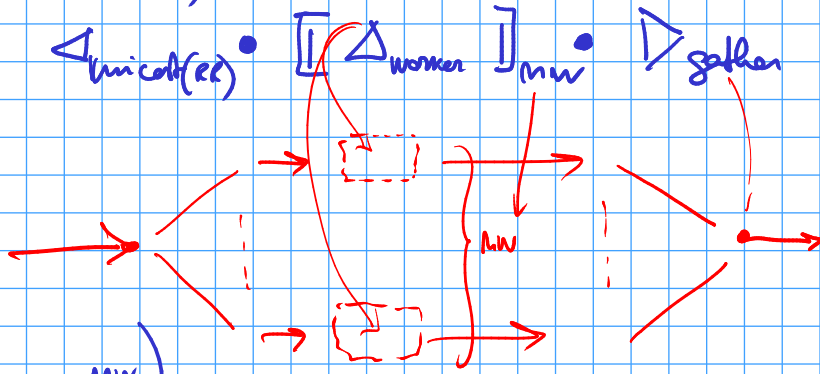


$$\Delta^{11} ::= ((f)) \mid ((g)) \mid \Delta \cdot \Delta \mid \Delta^{1m} \cdot \Delta^{m1} \mid \Delta^{1m} \cdot \Delta^m \cdot \Delta^{m1}$$

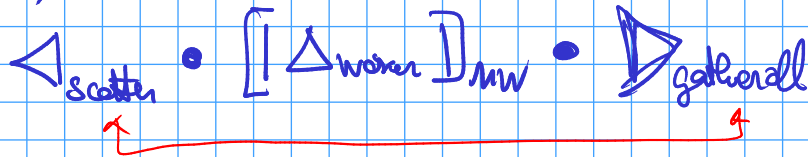


EMBARRASSINGLY PARALLEL

MAP (worker, mw) =



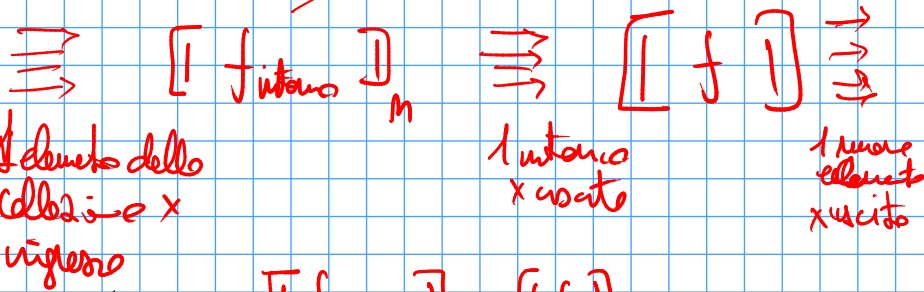
MAP (worker, mw) =



STENCIL (f, f intorno, mw)

Valore della collezione $x : y = f_{intorno}(x)$

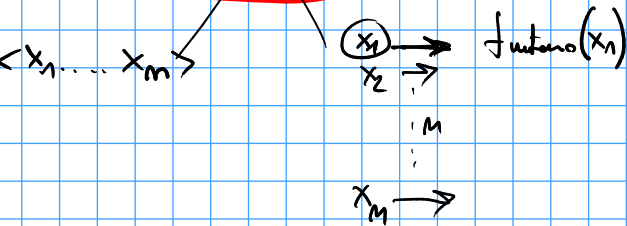
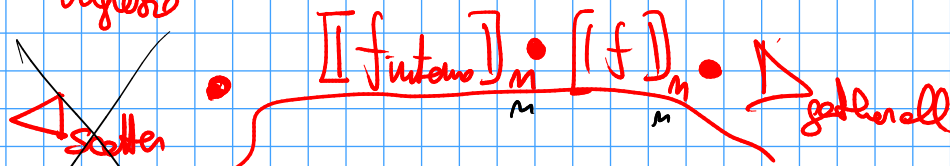
$f(y)$



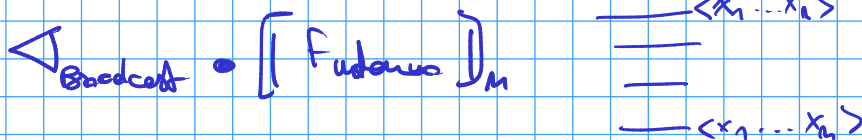
1 elemento della collezione x viene

1 istanza x viene

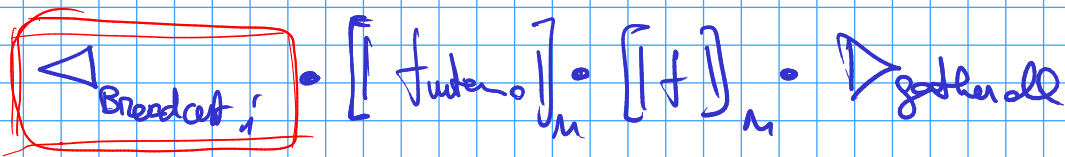
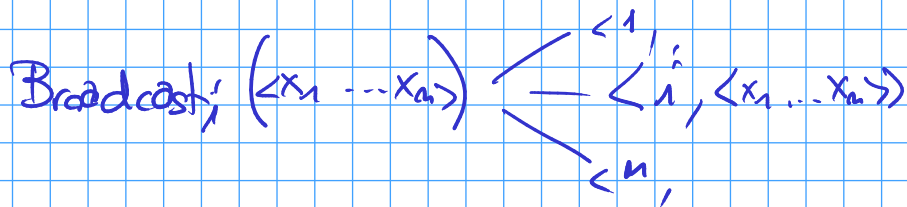
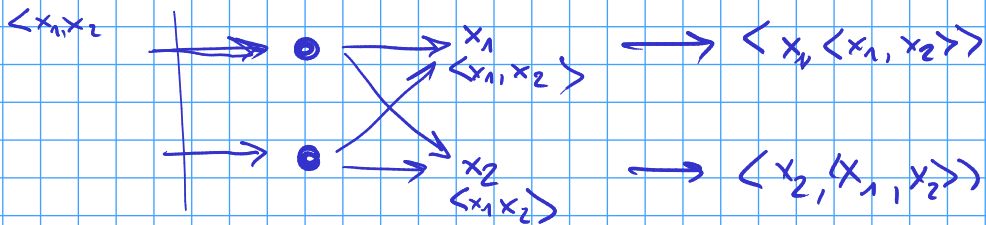
1 nuovo elemento x viene



$f_{intorno}(x_i) = \{x_i, x_{i-1}, x_{i+1}\}$



[Scatter, Broadcast]



Shared Memory Multicore

Broadcast_i: manages a queue worker
 Struct of int indice; float * x }

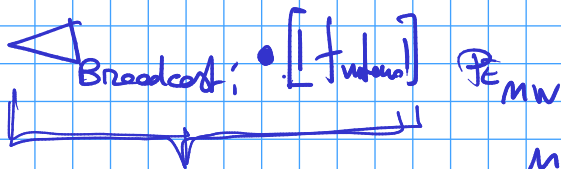
code "a mano" for: for_i
 Struct { float, float, float }
 i-2 i i+2

row/row

PE₀

PE₂

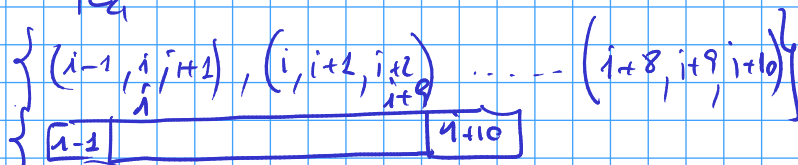
$$\frac{m}{MW} = 10$$



$$\underline{\underline{MW \ll m}}$$



PE₄

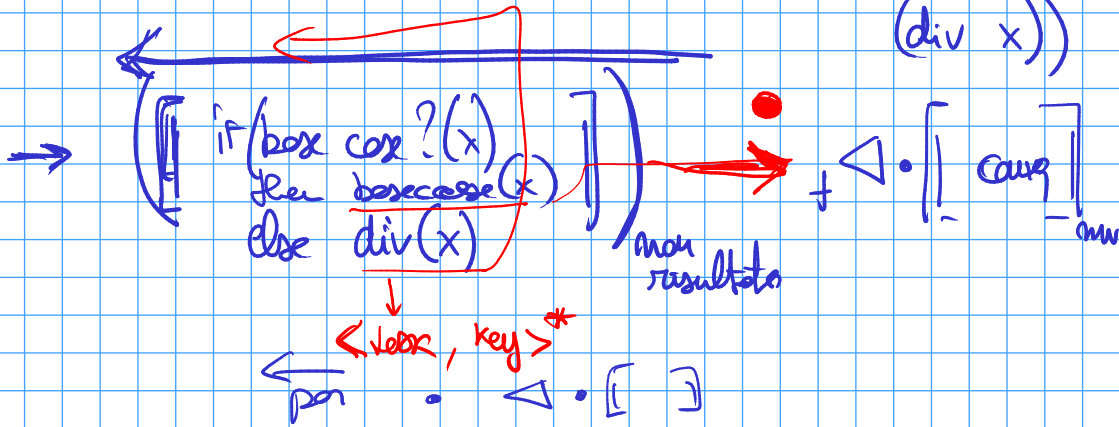


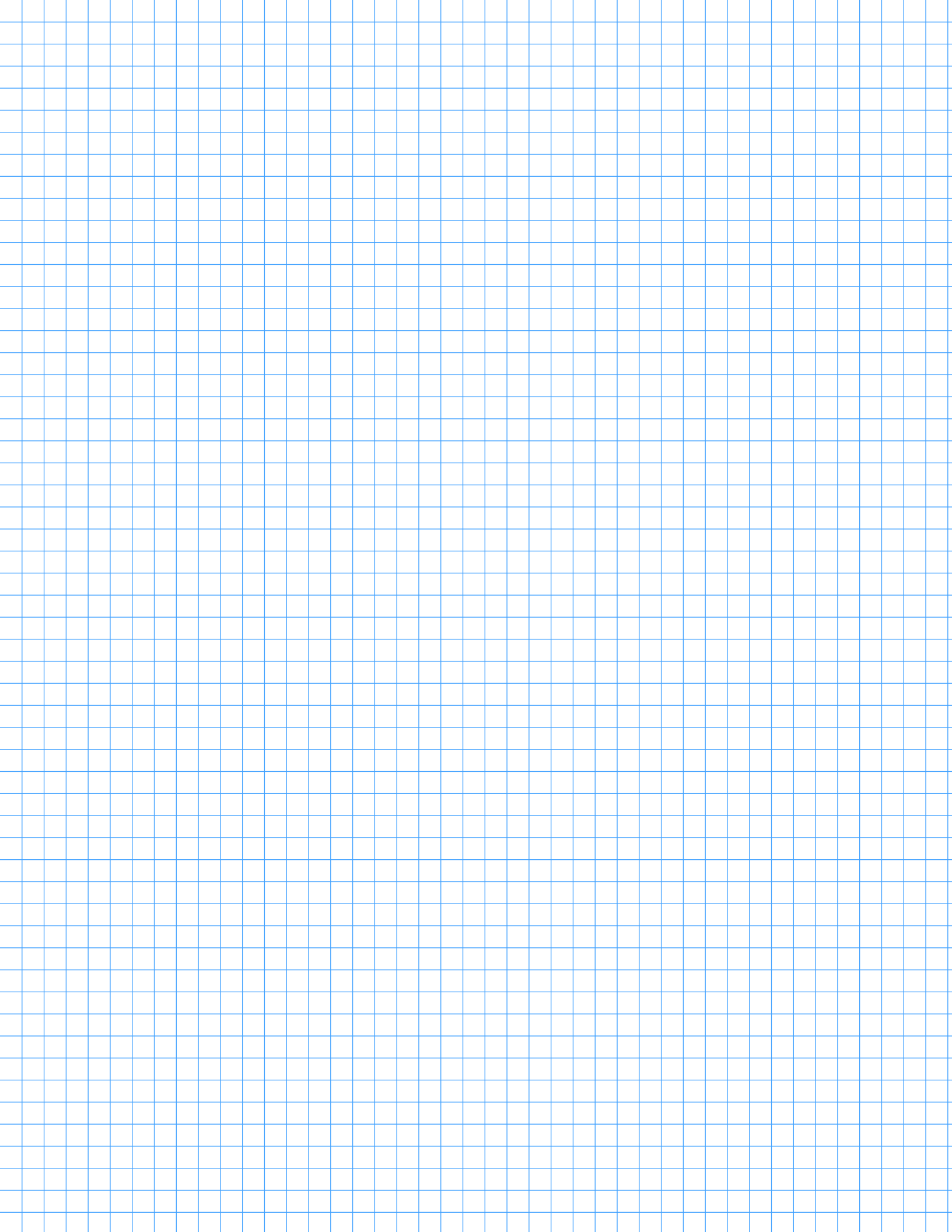
Divide & Conquer

div
conq

basecase?
basecase

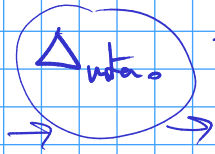
```
if (basecase?(x))
    then basecase(x);
else conq (apply to all
           sub
           (div x))
```





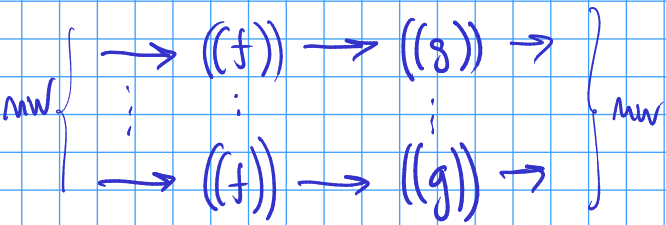
$$\left[\int_{\text{intorno}} \right]_{m \times n}$$

$((f_{\text{intorno}}))$

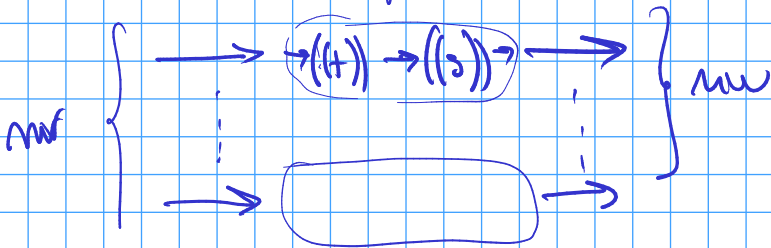


\Rightarrow building block
circulation

$$\left[((f)) \right]_{m \times n} \cdot \left[((g)) \right]_{m \times n} =$$

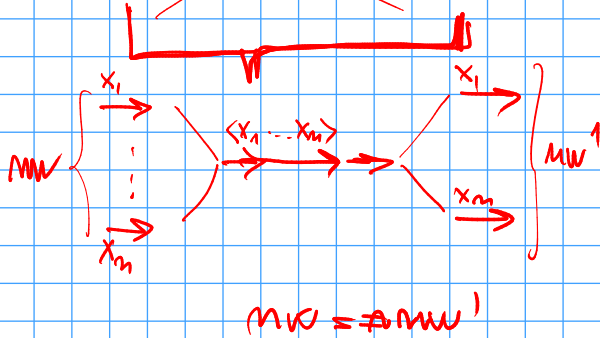


$$= \left[((f)) \cdot ((g)) \right]_{m \times n}$$

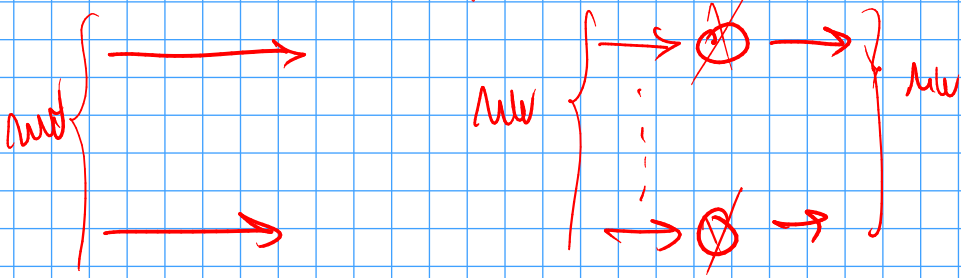


$$\text{map}(f) \circ \text{map}(g) = \text{map}(f \circ g)$$

$$\triangleleft_{\text{Satten}} \cdot \left[\left(\text{id}_{\text{Id}} \right) \right]_{\text{mw}} \cdot \triangleright_{\text{galkend}} \cdot \triangleleft_{\text{Satten}} \cdot \left[\left(g \right) \right]_{\text{mw}'} \cdot \triangleright_{\text{galkend}}$$



$$\left(\text{id}_{\text{Id}} \right) \quad \left[\left(\text{id}_{\text{Id}} \right) \right]_{\text{mw}}$$



$$\triangleleft_{\text{Satten}} \cdot \left[\left(f \right) \right]_{\text{mw}} \cdot \left[\left(g \right) \right]_{\text{mw}} \cdot \triangleright_{\text{galkend}}$$

$$\text{map}(f \circ g)$$